Appendix
MathLang: experience-driven development of a new mathematical language

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This document is an appendix to [KMW]. It is composed by the translation of the Foundations of Analysis’ first chapter into MathLang (Section A) and by E. Landau’s original text (Section B).

References


A Translation of the Foundations of Analysis’ first chapter

This section includes the full translation of the first chapter of E. Landau’s Foundations of Analysis [Lan51] in MathLang. The original translation is an XML file which we input by hand into the machine. An XSL transformation is then carried out to obtain a \LaTeX\ code which is then compiled into the rendering below. Our goal is to go further in the future and develop a software that transforms this rendering into a text which is closer to the English text originally written by the author. The following rendering is mainly experimental. It should not be seen as an outcome for MathLang.

Beginning of our MathLang translation of the first chapter of Foundations of Analysis [Lan51]

\[
\begin{array}{ll}
\vdash \forall x \forall y \forall z : N, x = y \land y = z \vdash x = z \\
\vdash x : N, y : N, x = y \vdash y = x \\
\vdash x : N, y : N, z : N, x = y, y = z \vdash x = z \\
\vdash \text{Ax1}() := 1 : \text{natural number} \\
\vdash \text{Ax1}() := 1 : N \\
\vdash \text{Ax1}() := 1 : N \\
\vdash x : N \vdash \text{Ax2}(x) := \exists y : y : \text{successor}(x) \\
\vdash \forall y : \text{Ax2}(x) \\
\vdash x : N, y : N, x = y \vdash \text{CoAx2}(x, y) := x' = y'
\end{array}
\]

Section 1: Axioms

\[1\]
\[2\]
\[3\]
\[4\]
\[5\]
\[6\]
\[7\]
\[8\]
\[9\]

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Section 2: Addition

Proof Theorem 1

\[ x : N, y : N, x \neq y \vdash x' \neq y' \]

Proof Theorem 3

Definition 1

Proof Theorem 4 part A \{a, b\}

Proof Theorem 4 part A I

Proof Theorem 4 part A II

Proof Theorem 4 part B

Proof Theorem 4 part B I

Proof Theorem 4 part B II
\[ x + y = z \]
\[ x : \mathbb{N}, y : \mathbb{N}, z : \mathbb{N} \]

\[ x + y = x + y \]
\[ (x + y) + 1 = (x + y) + 1 \]
\[ (x + y) + (y + z) = x + (y + z) \]
\[ x + y + z = x + y + z \]

\[ x : \mathbb{N}, y : \mathbb{N}, z : \mathbb{N} \]

\[ x + y = y + x \]
\[ (x + y) + y = y + (x + y) \]
\[ (x + y) + z = x + (y + z) \]
\[ x + (y + z) = (x + y) + z \]

\[ x : \mathbb{N}, y : \mathbb{N}, z : \mathbb{N} \]

\[ x + y = x + y \]
\[ (x + y) + 1 = (x + y) + 1 \]
\[ (x + y) + (y + z) = x + (y + z) \]
\[ x + y + z = x + y + z \]

\[ x : \mathbb{N}, y : \mathbb{N}, z : \mathbb{N} \]

\[ x + y = x + y \]
\[ (x + y) + 1 = (x + y) + 1 \]
\[ (x + y) + (y + z) = x + (y + z) \]
\[ x + y + z = x + y + z \]

\[ x : \mathbb{N}, y : \mathbb{N}, z : \mathbb{N} \]
\[\begin{align*}
(151), \ (\text{Def} + (\text{cm}) \triangleright x &= 1 \text{ or } \exists z : x = 1 + u \\
(152) &\triangleright \text{case1}(x, 1) \text{ or case2}(x, 1) \\
(153) &\triangleright 1 : \text{sn} \\
\end{align*}\]

**Proof Theorem 9 part B II** \(\{2.5, 7.1.2.2\}\)

\[\begin{align*}
y : \text{sn} \\
\text{case1}(x, y) &\triangleright y = y + 1 \\
\text{Def} &\triangleright y + 1 = x + 1 \\
(155), \ (156) &\triangleright y' = x + 1 \\
(157) &\triangleright \text{case3}(x, y') \\
\text{case2}(x, y) &\triangleright u : \text{sn} \\
\text{Def case2} &\triangleright x = y + u \\
\end{align*}\]

\[\begin{align*}
u = 1 \\
(159) \triangleright x &= y + 1 \\
(160) &\triangleright \text{case1}(x, y') \\
(161) &\triangleright \text{case3}(x, y') \\
\end{align*}\]

\[\begin{align*}
u \neq 1 \\
(163), \ (\text{Def} + (\text{cm})) &\triangleright u = 1 + w \\
(164) &\triangleright x = y + (1 + w) \\
\text{Th}(y, 1, w) &\triangleright x = y + (1 + w) \\
(166) &\triangleright x = y + w \\
(167) &\triangleright x = y' + w \\
(168) &\triangleright \text{case2}(x, y') \\
\end{align*}\]

\[\begin{align*}
u : \text{sn} \\
\text{Def case3} &\triangleright y = x + v \\
(170) &\triangleright x = y + v \\
(171) &\triangleright \text{case1}(x, y') \\
(172) &\triangleright y' = x + v' \\
\end{align*}\]

\(\text{Section 3: Ordering}\)

\[\begin{align*}
x : \text{sn}, \ y : \text{sn} &\triangleright x > y \iff \exists u : x = y + u \\
x : \text{sn}, \ y : \text{sn} &\triangleright x < y \iff \exists u : y = x + u \\
x : \text{sn}, \ y : \text{sn} &\triangleright \text{Th10}(x, y) \iff x = y \text{ or } x > y \text{ or } x < y \\
\end{align*}\]

**Proof Theorem 10** \(\{3.1\}\)

**Proof Theorem 11** \(\{3.2\}\)

\[\begin{align*}
x : \text{sn} \\
x > y &\triangleright \exists u : x = y + u \\
y < x &\triangleright \exists u : x = y + u \\
(183), \ (184) &\triangleright \forall x, \forall y : \text{sn} \triangleright \text{Th11}(x, y) \\
\end{align*}\]

**Proof Theorem 12** \(\{3.3\}\)

\[\begin{align*}
x : \text{sn} \\
x < y &\triangleright \exists u : y = x + v \\
y > x &\triangleright \exists u : y = x + v \\
(187), \ (188) &\triangleright \forall x, \forall y : \text{sn} \triangleright \text{Th12}(x, y) \\
\end{align*}\]

**Proof Theorem 13** \(\{3.4\}\)

\[\begin{align*}
x : \text{sn} \\
x = y &\triangleright \text{Th13}(x, y) \\
(193), \ (194) &\triangleright \forall x, \forall y : \text{sn} \triangleright \text{Th14}(x, y) \\
\end{align*}\]

**Proof Theorem 14** \(\{3.5\}\)

**Preliminary Remark** \(\{3.6\}\)
\( z : N \)
\( x : N \)
\( y : N \)

\[ z > y \]

\( \text{Th}19(x, y, z) \vdash x + z > y + z \)  
(241)

\( \text{Th}6(y, z) \vdash y + z = y + z \)  
(242)

\( \text{(Def Th}19\{_{222}\}) \vdash x + y > u + y \)  
(243)

\( \text{Th}6(u, y) \vdash u + y = y + u \)  
(244)

\( (242), (243), (244) \vdash y + z > y + u \)  
(245)

\( (241), (245) \vdash \text{Th}21(x, y, z, u) \)  
(246)

\( x : N, y : N, z : N, u : N, x > y, z > u \vdash \text{Th}22(x, y, z, u) \vdash x + z > y + u \)  
(247)

\( x : N, y : N, z : N, u : N, x > y, z > u \vdash \text{Th}22(x, y, z, u) \vdash x + z > y + u \)  
(248)

**Proof Theorem 22**  
(3.14)

\( x : N \)
\( y : N \)
\( z : N \)

\( x = y, z > u, \text{Th}19(z, u, x) \vdash x + z > y + u \)  
(249)

\( x > y, z > u, \text{Th}10(x, y, z) \vdash x + z > y + u \)  
(250)

\( x > y, z > u, \text{Th}21(x, y, z, u) \vdash x + z > y + u \)  
(251)

\( (249), (250), (251) \vdash \text{Th}22(x, y, z, u) \)  
(252)

\( x : N, y : N, z : N, u : N, x > y, z > u \vdash \text{Th}23(x, y, z, u) \vdash x + z > y + u \)  
(253)

**Proof Theorem 23**  
(3.15)

\( x : N \)
\( y : N \)
\( z : N \)

\( x = y, z = u \vdash x + z = y + u \)  
(254)

\( x > y, z > u \vdash \text{Th}22(x, y, z, u) \vdash x + z > y + u \)  
(255)

\( (254), (255), (256) \vdash \text{Th}23(x, y, z, u) \)  
(257)

\( x : N \vdash \text{Th}24(x) \vdash x \geq 1 \)  
(258)

**Proof Theorem 24**  
(3.16)

\( x : N \)
\( y : N \)

\( y > x \)

\( \text{(Def Th}21\{_{222}\}) \vdash y = x + u \)  
(259)

\( \text{Th}24(u) \vdash u \geq 1 \)  
(260)

\( (259), (260) \vdash \text{Th}25(x, y) \vdash y > x + 1 \)  
(261)

\( \text{Th}25(x, y) \vdash y > x + 1 \)  
(262)

\( x : N, y : N, y > x \vdash \text{Th}26(x, y) \vdash y \leq x \)  
(263)

**Proof Theorem 25**  
(3.17)

\( x : N \)
\( y : N \)

\( y > x \)

\( \text{Th}25(y, x) \vdash y > x + 1 \)  
(264)

**Proof Theorem 26**  
(3.18)

\( y : N \)

\( -y \leq x \)

\( \text{Th}25(y, x) \vdash y > x + 1 \)  
(265)

\( \text{Th}26(x, y) \vdash y \leq x \)  
(266)

**Proof Theorem 27**  
(3.19)

\( y : N \)

\( y > x \)

\( \text{Th}25(y, x) \vdash y > x + 1 \)  
(267)

**Proof Theorem 28**  
(3.20)

\( y : N \)

\( y > x \)

\( \text{Th}25(y, x) \vdash y > x + 1 \)  
(268)

**Proof Theorem 29**  
(3.21)

\( y : N \)

\( y > x \)

\( \text{Th}25(y, x) \vdash y > x + 1 \)  
(269)

**Proof Theorem 30**  
(3.22)
\[ V_{x,y,z}(x, y, z) \]

\[ \text{Def} \quad (x, y) \mapsto (x, y, z) \]

\[ x + 1 = y \]

\[ (4.2) \quad x + 1 = y \]

\[ (327), (328) \quad x + 1 = y \]

\[ (329) \quad x + 1 = z \]

\[ (330) \quad x + 1 = z \]

\[ \text{Proof Theorem 30 part 1} \]

\[ (4.5.2) \]

\[ x : 2N \]

\[ \text{Th30}(x, y, z) \]

\[ (332) \]

\[ (333) \]

\[ (334) \]

\[ (335) \]

\[ (336) \]

\[ (337) \]

\[ (338) \]

\[ \text{Ax5}(2N, (331), (339)) \]

\[ \text{Proof Theorem 30} \]

\[ (4.5) \]

\[ \text{Proof Theorem 31} \]

\[ (4.6.1) \]

\[ x : 1N \]

\[ (39) : 51N \]

\[ \text{Th30}(x, y, z) \]

\[ (348) \]

\[ (349) \]

\[ (350) \]

\[ (351) \]

\[ (352) \]

\[ (353) \]

\[ (354) \]

\[ (355) \]

\[ \text{Ax5}(2N, (348), (356)) \]

\[ \text{Proof Theorem 31} \]

\[ (4.6) \]

\[ z : 3N \]

\[ \text{Th31}(x, y, z) \]

\[ (360) \]

\[ (361) \]

\[ (362) \]

\[ (363) \]

\[ (364) \]

\[ \text{Proof Theorem 31 part 1} \]

\[ (4.6.2) \]

\[ x : 1N \]

\[ (39) : 51N \]

\[ \text{Th31}(x, y, z) \]

\[ (365) \]

\[ (366) \]

\[ (367) \]

\[ (368) \]

\[ (369) \]

\[ (370) \]

\[ (371) \]

\[ \text{Ax5}(2N, (364), (372)) \]

\[ \text{Proof Theorem 32} \]

\[ (4.7) \]

\[ x : 1N \]

\[ y : 1N \]

\[ \text{Th32}(x, y, z) \]

\[ (374) \]

\[ (375) \]

\[ (376) \]

\[ (377) \]

\[ \text{Proof Theorem 32 case 1} \]

\[ (4.7.1) \]
\[\text{Def } \triangleright \vdash x = y + u\]  
\[\begin{align*}  
\text{1978} & : x \geq y, z 
\text{1981} & : x = (y + u), z 
\end{align*}\]  
\[\text{Th30}(z, y, u) \vdash (y + u), z = y, z + u, z\]

\[\text{Def } \triangleright \vdash y, z + u, z \geq y, z\]  
\[\text{(1979), (1980), (1981) } \vdash z, y \geq y, z\]  
\[x = y\]

\[\vdash x \geq y\]  
\[\vdash \text{Proof Theorem 32 case 2}\]  
\[\begin{align*}  
\text{(384)} & : x \geq y, z = x, z 
\text{(385)} & : x < y, z = x, z 
\text{(386)} & : x = y, z = x, z 
\end{align*}\]

\[\vdash \text{Proof Theorem 32 case 3}\]  
\[\begin{align*}  
\text{(387)} & : x \geq y, z = y, z 
\text{(388)} & : x = y, z = y, z 
\end{align*}\]

\[\vdash \text{Proof Theorem 33}\]  
\[\begin{align*}  
\text{(390)} & : x = y, z = y, z 
\end{align*}\]

\[\vdash \text{Proof Theorem 34}\]  
\[\begin{align*}  
\text{(391)} & : x = y, z = y, z 
\end{align*}\]

\[\vdash \text{Proof Theorem 35}\]  
\[\begin{align*}  
\text{(392)} & : x = y, z = y, z 
\end{align*}\]

\[\vdash \text{Proof Theorem 36}\]  
\[\begin{align*}  
\text{(393)} & : x = y, z = y, z 
\end{align*}\]

End of our MathLang translation of the first chapter of
\textit{Foundations of Analysis} [Lan51]

\section*{B Foundations of Analysis' first chapter}

This section of the appendix contains the original text of the first chapter (translated from German to English by F. Steinhardt) of E. Landau’s \textit{Foundations of Analysis} [Lan51, Lan30].

Beginning of the original first chapter of \textit{Foundations of Analysis} [Lan51]
B.1 Natural Numbers

B.1.1 Axioms

We assume the following to be given:
A set (i.e. totality) of objects called natural numbers, possessing the properties - called axioms- to be listed below.

Before formulating the axioms we make some remarks about the symbols = and \( \neq \) which will be used.

Unless otherwise specified, small italic letters will stand for natural numbers throughout this book.

If \( x \) is given and \( y \) is given, then either \( x \) and \( y \) are the same number; this may be written

\[ x = y \]

(= to be read “equals”); or \( x \) and \( y \) are not the same number; this may be written

\[ x \neq y \]

(\( \neq \) to be read “is not equal to”).

Accordingly, the following are true on purely logical grounds:

\[ x = x \text{ for every } x \tag{1} \]

\[ \text{If } x = y \text{ then } y = x \tag{2} \]

\[ \text{If } x = y, y = z \text{ then } x = z \tag{3} \]

Thus a statement such as

\[ a = b = c = d, \]

which on the face of it means merely that

\[ a = b, b = c, c = d, \]

contains the additional information that, say,

\[ a = c, a = d, b = d. \]

(Similarly in the later chapters.)

Now, we assume that the set of all natural numbers has the following properties:

**Axiom 1** 1 is a natural number.

That is, our set is not empty; it contains an object called 1 (read “one”).

**Axiom 2** For each \( x \) there exists exactly one natural number, called the successor of \( x \), which will be denoted by \( x' \).

In the case of complicated natural numbers \( x \), we will enclose in parentheses the number whose successor is to be written down, since otherwise ambiguities might arise. We will do the same throughout this book, in case of \( x + y, xy, x - y, -x, x^2 \), etc.

Thus, if

\[ x = y \]

then

\[ x' = y'. \]

**Axiom 3** We always have

\[ x' \neq 1. \]

That is, there exists no number whose successor is 1.

**Axiom 4** If

\[ x' = y' \]

then

\[ x = y. \]

That is, for any given number there exists either no number or exactly one number whose successor is the given number.

**Axiom 5 (Axiom of Induction)** Let there be given a set \( \mathbb{N} \) of natural numbers, with the following properties:

1) 1 belongs to \( \mathbb{N} \)

2) If \( x \) belongs to \( \mathbb{N} \) then so does \( x' \)

Then \( \mathbb{N} \) contains all the natural numbers.

B.1.2 Addition

**Theorem 1** If

\[ x \neq y \]

then

\[ x' \neq y'. \]

**Proof** Otherwise, we would have

\[ x' = y' \]

and hence, by Axiom 4,

\[ x = y. \]

**Theorem 2**

\[ x' \neq x. \]

**Proof** Let \( \mathbb{N} \) be the set of all \( x \) for which this holds true.

1) By Axiom 1 and Axiom 3,

\[ 1' \neq 1; \]

therefore 1 belongs to \( \mathbb{N} \).

2) If \( x \) belongs to \( \mathbb{N} \), then

\[ x' \neq x, \]

and hence by Theorem 1,

\[ (x')' \neq x', \]

so that \( x' \) belongs to \( \mathbb{N} \).

By Axiom 5, \( \mathbb{N} \) therefore contains all the natural numbers, i.e. we have for each \( x \) that

\[ x' \neq x. \]

**Theorem 3** If

\[ x \neq 1, \]

then there exists one (hence, by Axiom 4, exactly one) \( u \) such that

\[ x = u'. \]
**Proof** Let \( \mathbb{N} \) be the set consisting of the number 1 and of all those \( x \) for which there exists such a \( u \). (For any such \( x \), we have of necessity that 
\[
x \neq 1
\]
by Axiom 3.)

I) \( 1 \) belongs to \( \mathbb{N} \).

II) If \( x \) belongs to \( \mathbb{N} \), then, with \( u \) denoting the number \( x \), we have 
\[
x' = u',
\]
so that \( x' \) belongs to \( \mathbb{N} \).
By Axiom 5, \( \mathbb{N} \) therefore contains all the natural numbers; thus for each 
\[
x \neq 1
\]
there exists a \( u \) such that 
\[
x = u'.
\]

**Theorem 4**, and at the same time **Definition 1** To every pair of numbers \( x, y \), we may assign in exactly one way a natural number, called \( x + y \) (+ to be read “plus”), such that 
\[
x + 1 = x' \quad \text{for every } x \quad (4)
\]
\[
x + y' = (x + y)' \quad \text{for every } x \text{ and every } y \quad (5)
\]
\( x + y \) is called the sum of \( x \) and \( y \), or the number obtained by addition of \( y \) to \( x \).

**Proof**

A) First we will show that for each fixed \( x \) there is at most one possibility of defining \( x + y \) for all \( y \) in such a way that 
\[
x + 1 = x'
\]
and 
\[
x + y' = (x + y)' \quad \text{for every } y.
\]
Let \( a_y \) and \( b_y \) be defined for all \( y \) and be such that 
\[
a_1 = x', \quad b_1 = x';
\]
\[
a_y = (a_y)' \quad \text{for every } y.
\]
Let \( \mathbb{N} \) be the set of all \( y \) for which 
\[
a_y = b_y.
\]

I) 
\[
a_1 = x' = b_1;
\]

hence \( 1 \) belongs to \( \mathbb{N} \).

II) If \( y \) belongs to \( \mathbb{N} \), then 
\[
a_y = b_y;
\]
hence by Axiom 2, 
\[
(a_y)' = (b_y)',
\]
therefore 
\[
a_y' = (a_y)' = (b_y)' = b_y',
\]
so that \( y' \) belongs to \( \mathbb{N} \).
Hence \( \mathbb{N} \) is the set of all natural numbers; i.e. for every \( y \) we have 
\[
a_y = b_y.
\]

B) Now we will show that for each \( x \) it is actually possible to define \( x + y \) for all \( y \) in such a way that 
\[
x + 1 = x'
\]
and 
\[
x + y' = (x + y)',
\]
Let \( \mathbb{N} \) be the set of all \( x \) for which this is possible (in exactly one way, by A)

I) For 
\[
x = 1,
\]
the number 
\[
x + y = y'
\]
is as required, since 
\[
x + 1 = 1' = x',
\]
\[
x + y' = (y')' = (x + y)'.
\]
Hence \( 1 \) belongs to \( \mathbb{N} \).

II) Let \( x \) belong to \( \mathbb{N} \), so that there exists an \( x + y \) for all \( y \). Then the number 
\[
x' + y = (x + y)'
\]
is the required number for \( x' \), since 
\[
x' + 1 = (x + 1)' = (x')',
\]
and 
\[
x' + y' = (x + y)' = [(x + y)']' = (x + y)'.
\]
Hence \( x' \) belongs to \( \mathbb{N} \).
Therefore \( \mathbb{N} \) contains all \( x \).

**Theorem 5 (Associative Law of Addition)**
\[
(x + y) + z = x + (y + z).
\]

**Proof** Fix \( x \) and \( y \), and denote by \( \mathbb{N} \) the set of all \( z \) for which the assertion of the theorem holds.

I) 
\[
(x + y) + 1 = (x + y)' = x + y' = x + (y + 1);
\]

thus \( 1 \) belongs to \( \mathbb{N} \).

II) Let \( z \) belong to \( \mathbb{N} \). Then 
\[
(x + y) + z = x + (y + z),
\]
hence 
\[
(x + y) + z' = [(x + y) + z]' = (x + (y + z))' = x + (y + z)'.
\]
so that \( z' \) belongs to \( \mathbb{N} \).
Therefore the assertion holds for all \( z \).

**Theorem 6 (Commutative Law of Addition)**
\[
x + y = y + x.
\]

**Proof** Fix \( y \), and \( \mathbb{N} \) be the set of all \( x \) for which the assertion holds.

I) We have 
\[
y + 1 = y',
\]
and furthermore, by the construction in the proof of Theorem 4, 
\[
1 + y = y',
\]
so that 
\[
1 + y = y + 1
\]
and \( 1 \) belongs to \( \mathbb{N} \).
II) If \( x \) belongs to \( \mathfrak{M} \), then
\[
x + y = y + x,
\]
Therefore
\[
(x + y)' = (y + x)' = y + x'.
\]
By the construction in the proof of Theorem 4, we have
\[
x' + y = (x + y)',
\]

1) hence
\[
x' + y = y + x',
\]
so that \( x' \) belongs to \( \mathfrak{M} \).
The assertion therefore holds for all \( x \).

**Theorem 7**
\( y \neq x + y \).

**Proof** Fix \( x \), and let \( \mathfrak{M} \) be the set of all \( y \) for which the assertion holds.

1) \( 1 \neq x' \),
\[
1 \neq x + 1;
\]
\( 1 \) belongs to \( \mathfrak{M} \).

II) If \( y \) belongs to \( \mathfrak{M} \), then
\[
y \neq x + y,
\]
hence
\[
y' \neq (x + y)',
\]
y' \neq x + y',
so that \( y' \) belongs to \( \mathfrak{M} \).
Theorem 8 follows.

**Theorem 8** If
\( y \neq z \)
Then
\( x + y \neq x + z \).

**Proof** Consider a fixed \( y \) and a fixed \( z \) such that
\( y \neq z \),
and let \( \mathfrak{M} \) be the set of all \( x \) for which
\[
x + y \neq x + z,
\]

1) \( y' \neq x' \),
\[
1 + y \neq 1 + z;
\]
hence \( 1 \) belongs to \( \mathfrak{M} \).

II) If \( x \) belongs to \( \mathfrak{M} \), then
\[
x + y \neq x + z,
\]
hence
\[
(x + y)' \neq (x + z)',
\]
\( x' + y \neq x' + z \),
so that \( x' \) belongs to \( \mathfrak{M} \).
The assertion therefore holds always.

**Theorem 9** For given \( x \) and \( y \), exactly one of the following must be the case:

1) \( x = y \).

2) There exists a \( u \) (exactly one, by Theorem 8) such that
\[ x = y + u \]

3) There exists a \( v \) (exactly one, by Theorem 8) such that
\[ y = x + v \]

**Proof**
A) By Theorem 7, cases 1) and 2) are incompatible. Similarly, 1) and 3) also follows from Theorem 7; for otherwise, we would have
\[
x = y + u = (x + v) + u = x + (v + u) = (v + u) + x.
\]
Therefore we can have at most one of the cases 1), 2) and 3).

B) Let \( x \) be fixed, and let \( \mathfrak{M} \) be the set of all \( y \)
for which one (hence by A), exactly one) of the cases 1), 2) and 3) obtains.

I) For \( y = 1 \), we have by Theorem 3 that either
\[
x = 1 = y
\]

\( \text{(case 1)} \)
or
\[
x = u' = 1 + u = y + u
\]

\( \text{(case 2)} \).

Hence \( 1 \) belongs to \( \mathfrak{M} \).

II) Let \( y \) belong to \( \mathfrak{M} \). Then either (case 1) for \( y \)
\[
x = y,
\]
hence
\[
y' = y + 1 = x + 1
\]

\( \text{(case 3) for } y' \);
or (case 2) for \( y \)
\[
x = y + u,
\]
hence if
\[
u = 1,
\]
then
\[
x = y + 1 = y'
\]

\( \text{(case 1) for } y' \);
but if
\[
u \neq 1,
\]
then, by Theorem 3,
\[
u = w' = 1 + w,
\]
\[
x = y + (1 + w) = (y + 1) + w = y' + w
\]

\( \text{(case 2) for } y' \);
or (case 3) for \( y \)
\[
y = x + v,
\]
hence
\[
y' = (x + v)' = x + v'
\]

\( \text{(case 3) for } y' \).

In any case, \( y' \) belongs to \( \mathfrak{M} \).

Therefore we always have one of the cases 1), 2) and 3).
B.1.3 Ordering

Definition 2 If
\[ x = y + u \]
then
\[ x > y. \]
(> to be read “is greater than.”)

Definition 3 If
\[ y = x + v \]
then
\[ x < y. \]
(< to be read “is less than.”)

Theorem 10 For any given \( x, y \), we have exactly one of the cases
\[ x = y, x > y, x < y. \]
Proof Theorem 9, Definition 2 and Definition 3.

Theorem 11 If
\[ x > y \]
then
\[ y < x. \]
Proof Each of these means that
\[ x = y + u \]
for some suitable \( u \).

Theorem 12 If
\[ x < y \]
then
\[ y > x. \]
Proof Each of these means that
\[ y = x + v \]
for some suitable \( v \).

Definition 4
\[ x \geq y \]
means
\[ x > y \text{ or } x = y. \]
(\( \geq \) to be read “is greater than or equal to.”)

Definition 5
\[ x \leq y \]
means
\[ x < y \text{ or } x = y. \]
(\( \leq \) to be read “is less than or equal to.”)

Theorem 13 If
\[ x \geq y \]
then
\[ y \leq x. \]
Proof Theorem 11.

Theorem 14 If
\[ x \leq y \]
then
\[ y \geq x. \]
Proof
\[ x + y = x + y \]

Proof Theorem 12.

Theorem 15 (Transitivity of Ordering) If
\[ x < y, y < z, \]
then
\[ x < z. \]

Preliminary Remark Thus if
\[ x > y, y > z, \]
then
\[ x > z, \]
since
\[ z < y, y < x, z < x; \]
but in what follows I will not even bother to write down such statements, which are obtained trivially by simply reading the formulas backwards.

Proof With suitable \( v, w \), we have
\[ y = x + v, z = y + w, \]

\[ z = (x + v) + w = x + (v + w), x < z. \]

Theorem 16 If
\[ x \leq y, y < z, \text{ or } x < y, y \leq z, \]
then
\[ x < z. \]

Proof Obvious if an equality sign holds in the hypothesis; otherwise, Theorem 15 does it.

Theorem 17 If
\[ x \leq y, y \leq z, \]
then
\[ x \leq z. \]

Proof Obvious if two equality signs hold in the hypothesis; otherwise, Theorem 16 does it.

A notation such as
\[ a < b \leq c < d \]
is justified on the basis of Theorem 15 and 17. While its immediate meaning is
\[ a < b, b \leq c, c < d, \]
it also implies, according to these theorems, that, say
\[ a < c, a < d, b < d. \]
(Similarly in the later chapters.)

Theorem 18
\[ x + y > x. \]

Proof
\[ x + y = x + y \]
Theorem 19 If
\[ x > y, \text{ or } x = y, \text{ or } x < y, \]
then
\[ x + z > y + z, \text{ or } x + z = y + z, \text{ or } x + z < y + z, \]
respectively.

Proof
1) If \( x > y \)
then
\[ x = y + u, \]
\[ x + z = (y + u) + z = (u + y) + z = u + (y + z) = \]
\[ (y + z) + u, \]
\[ x + z > y + z. \]

2) If \( x = y \)
then clearly
\[ x + z = y + z. \]

3) If \( x < y \)
then
\[ y > x, \]
and hence, by 1),
\[ y + z > x + z, \]
\[ x + z < y + z. \]

Theorem 20 If
\[ x + z > y + z, \text{ or } x + z = y + z, \text{ or } x + z < y + z, \]
then
\[ x > y, \text{ or } x = y, \text{ or } x < y, \]
respectively.

Proof Follows from Theorem 19, since the three cases are, in both instances, mutually exclusive and exhaust all possibilities.

Theorem 21 If
\[ x > y, z > u, \]
then
\[ x + z > y + u. \]

Proof By Theorem 19, we have
\[ x + z > y + z \]
and
\[ y + z = z + y > u + y = y + u \]
hence
\[ x + z > y + u. \]

Theorem 22 If
\[ x \geq y, z > u \text{ or } x > y, z \geq u, \]
then
\[ x + z \geq y + u. \]

Proof Follows Theorem 19 if an equality sign holds in the hypothesis, otherwise from Theorem 21.

Theorem 23 If
\[ x > y, z > u, \]
then
\[ x + z > y + u. \]

Proof Obvious if two equality signs hold in the hypothesis; otherwise Theorem 22 does it.

Theorem 24
\[ x \geq 1. \]

Proof Either
\[ x = 1 \]
or
\[ x = u' = u + 1 > 1. \]

Theorem 25 If
\[ y > x \]
then
\[ y \geq x + 1. \]

Proof
\[ y = x + u, \]
\[ u \geq 1, \]
hence
\[ y \geq x + 1. \]

Theorem 26 If
\[ y < x + 1 \]
then
\[ y \leq x. \]

Proof Otherwise we would have
\[ y > x \]
and therefore, by Theorem 25,
\[ y \geq x + 1. \]

Theorem 27 In every non-empty set of natural numbers there is at least one (i.e., one which is less than any other number of the set).

Proof Let \( \mathbb{N} \) be the given set, and let \( \mathbb{N} \) be the set of all \( x \) which are \( \leq \) every number \( n \).

By Theorem 24, the set \( \mathbb{N} \) contains the number 1. Not every \( x \) belongs to \( \mathbb{N} \); in fact for each \( y \) of \( \mathbb{N} \) the number \( y + 1 \) does not belong to \( \mathbb{N} \), since
\[ y + 1 > y. \]

Therefore there is an \( m \) in \( \mathbb{N} \) such that \( m + 1 \) does not belong to \( \mathbb{N} \); for otherwise, every natural number would have to belong to \( \mathbb{N} \), by Axiom 5.

Of this \( m \) I now assert that it is \( \leq \) every \( n \) of \( \mathbb{N} \), and that it belongs to \( \mathbb{N} \). The former we already know.

The latter is established by an indirect argument, as follows: If \( m \) did not belong to \( \mathbb{N} \), then for each \( n \) of \( \mathbb{N} \) we would have
\[ m < n, \]
hence, by Theorem 25,
\[ m + 1 \leq n; \]
thus \( m + 1 \) would belong to \( \mathbb{N} \), contradicting the statement above by which \( m \) was introduced.
B.1.4 Multiplication

Theorem 28 and at the same time Definition 6. To every pair of numbers \(x, y\), we may assign in exactly one way a natural number, called \(x \cdot y\) (to be read “times”; however, the dot is usually omitted), such that

\[
x \cdot 1 = x \quad \text{for every } x,
\]

\[
x \cdot y' = (x \cdot y) + x \quad \text{for every } x \text{ and } y.
\]

\(x \cdot y\) is called the product of \(x\) and \(y\), or the number obtained from multiplication of \(x\) by \(y\).

Proof (mutatis mutandis, word for word the same as that of Theorem 4)

A) We will first show that for each fixed \(x\) there is at most one possibility of defining \(x y\) for all \(y\) in such a way that

\[
x \cdot 1 = x'
\]

and

\[
x y' = xy + x \quad \text{for every } y.
\]

Let \(a_y\) and \(b_y\) be defined for all \(y\) and be such that

\[
a_1 = x, \quad b_1 = x,
\]

\[
a_y' = a_y + x, \quad b_y = b_y + x \quad \text{for every } y.
\]

Let \(\mathbb{N}\) be the set of all \(y\) for which

\[
a_y = b_y.
\]

1) \(a_1 = x = b_1;\)

hence \(1\) belongs to \(\mathbb{N}\).

11) If \(y\) belongs to \(\mathbb{N}\), then

\[
a_y = b_y,
\]

hence,

\[
a_y' = a_y + x = b_y + x = b_y',
\]

so that \(y\) belongs to \(\mathbb{N}\).

Hence \(\mathbb{N}\) is the set of all natural numbers; i.e. for every \(y\) we have

\[
a_y = b_y.
\]

B) Now we will show that for each \(x\) it is actually possible to define \(xy\) for all \(y\) in such a way that

\[
x \cdot 1 = x
\]

and

\[
x y' = xy + x \quad \text{for every } y.
\]

Let \(\mathbb{N}\) be the set of all \(x\) for which this is possible (in exactly one way, by A))

1) For \(x = 1,\)

the number

\[
xy = y
\]

is as required, since

\[
x \cdot 1 = 1 = x,
\]

\[
x y' = y' = y + 1 = xy + x.
\]

Hence \(1\) belongs to \(\mathbb{N}\).

II) Let \(x\) belong to \(\mathbb{N}\), so that there exists an \(xy\) for all \(y\). Then the number

\[
x' y = xy + y
\]

is the required number for \(x\), since

\[
x' \cdot 1 = x \cdot 1 + 1 = x + 1 = x'
\]

and

\[
x' y' = xy' + y' = (xy + x) + y' =
\]

\[
xy + (x + y') = xy + (x' + y) =
\]

\[
xy + (y + x') = (xy + y) + x' = x'y + x'.
\]

Hence \(x\) belongs to \(\mathbb{N}\).

Therefore \(\mathbb{N}\) contains all \(x\).

Theorem 29 (Commutative Law of Multiplication)

\[xy = yx.\]

Proof Fix \(y\), and let \(\mathbb{N}\) be the set of all \(x\) for which the assertion holds.

1) We have

\[
y \cdot 1 = y,
\]

and furthermore, by the construction in the proof of Theorem 28,

\[
1 \cdot y = y,
\]

hence

\[
1 \cdot y = y \cdot 1,
\]

so that \(1\) belongs to \(\mathbb{N}\).

11) If \(x\) belongs to \(\mathbb{N}\), then

\[
x y = yx,
\]

hence

\[
x y + y = yx + y = yx'.
\]

By the construction in the proof of Theorem 28, we have

\[
x' y = xy + y,
\]

hence

\[
x' y = yx',
\]

so that \(x'\) belongs to \(\mathbb{N}\).

The assertion therefore holds for all \(x\).

Theorem 30 (Distributive Law)

\[x(y + z) = xy + xz.\]

Preliminary Remark The formula

\[(y + z)x = yz + zx\]

which results from Theorem 30 and Theorem 29, and similar analogues later on, need not be specifically formulated as theorems, nor even be set down.

Proof Fix \(x\) and \(y\), and let \(\mathbb{N}\) be the set of all \(z\) for which the assertion holds true.

1) \(x(y + z) = xy + z = xy + x \cdot 1;\)

1 belongs to \(\mathbb{N}\).

11) If \(z\) belongs to \(\mathbb{N}\), then

\[
x(y + z) = xy + xz,
\]

hence

\[
x(y + z') = x((y + z')) = x(y + z) + x =
\]

\[
(xy + z) + x = xy + (xz + x) = xy + xz',
\]

so that \(z'\) belongs to \(\mathbb{N}\).

Therefore, the assertion always holds.
Theorem 31 (Association Law of Multiplication)

\[(xy)z = x(yz)\].

**Proof** Fix \(x\) and \(y\), and let \(\mathbb{N}\) be the set of all \(z\) for which the assertion holds true.

1) \((xy) \cdot 1 = xy = x(y \cdot 1)\);

hence \(1\) belongs to \(\mathbb{N}\).

II) Let \(z\) belong to \(\mathbb{N}\). Then

\[(xy)z = x(yz),\]

and therefore, using Theorem 30,

\[(xy)z' = (xy)z + xy = x(yz' + y) = x(yz + y) = x(yz'),\]

so that \(z'\) belongs to \(\mathbb{N}\).

Therefore \(\mathbb{N}\) contains all natural numbers.

Theorem 32 If

\[x > y, \text{ or } x = y, \text{ or } x < y,\]

then

\[xz > yz, \text{ or } xy = yz, \text{ or } xz < yz, \text{ respectively.}\]

**Proof**

1) If \(x > y\)

then \(x = y + u,\)

\[xz = (y + u)z = yz + uz > yz.\]

2) If \(x = y\)

then clearly \(xz = yz.\)

3) If \(x < y\)

then \(y > x,\)

hence by 1),

\[yz > xz,\]

\[xz < yz.\]

Theorem 33 If

\[xz > yz, \text{ or } xz = yz, \text{ or } xz < yz,\]

then

\[x > y, \text{ or } x = y, \text{ or } x < y, \text{ respectively.}\]

**Proof** Follows from Theorem 32, since the three cases are, in both instances, mutually exclusive and exhaust all possibilities.

Theorem 34 If

\[x > y, z > u,\]

then

\[xz > yu.\]

**Proof** By Theorem 32, we have

\[xz > yz\]

and

\[yz = zy > uy = yu,\]

hence

\[xz > yu.\]

Theorem 35 If

\[x \geq y, z > u \text{ or } x > y, z \geq u,\]

then

\[xz > yu.\]

**Proof** Follows from Theorem 32 if an equality sign holds in the hypothesis; otherwise from Theorem 34.

Theorem 36 If

\[x \geq y, z \geq u,\]

then

\[xz \geq yu.\]

**Proof** Obvious if two equality signs hold in the hypothesis; otherwise Theorem 35 does it.